
Data Compression

LZ77

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Outline

- Introduction – Principle of dictionary methods
- LZ77 Sliding window
- Examples
- Optimization
- Performance comparison
- Applications/Patents

Principle of dictionary methods

- Compressing multiple strings can be more efficient than compressing single symbols only (e.g. Huffman encoding).
- Strings of symbols are added to a dictionary. Later occurrences are referenced.
- Static dictionary: Entries are predefined and constant according to the application of the text
- Adaptive dictionary: Entries are taken from the text itself and created on-the-fly

LZ77

- First paper by Ziv and Lempel in 1977 about lossless compression with an adaptive dictionary.
- Goes through the text in a **sliding window** consisting of a *search buffer* and a *look ahead buffer*.

Search buffer Look-ahead buffer

...this is a text that is being read through the window...

- The search buffer is used as dictionary
 - Sizes of these buffers are parameters of the implementation. Assumption: Patterns in text occur within range of the search buffer.
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LZ77 – Example (Encoding)

Encoding of the string:

abracadabrad

output tuple: (offset, length, symbol)

	7	6	5	4	3	2	1		output	
								a b r a c	ada...	(0,0,a)
						a	b r a c a	dab...	(0,0,b)	
					a	b r a c a d		abr...	(0,0,r)	
				a	b r a c a d a			bra...	(3,1,c)	
		a	b r a c a d a b r					ad	(2,1,d)	
...ac	a	b r a c a d a b r a d							(7,4,d)	
a	d a b r a d									

Search buffer	Look-ahead buffer
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12 characters compressed into 6 tuples

Compression rate: $(12 \cdot 8) / (6 \cdot (5 + 2 + 3)) = 96 / 60 = 1,6 = 60\%$.

Size of output

- Size for each output tuple (offset, length, symbol) when using fixed-length storage:

$$\lceil \log_2 S \rceil + \lceil \log_2 (S + L) \rceil + \lceil \log_2 A \rceil$$

where S is the length of the search buffer, L the length of the look ahead window, A the size of the alphabet.

- Why $S+L$ and not only S ? See next slide.
- Worst case if no symbol repeats in the search buffer:

Blow up of $n(\lceil \log_2 S \rceil + \lceil \log_2 (S + L) \rceil + \lceil \log_2 A \rceil)$
instead of $n \lceil \log_2 A \rceil$

Encoding reaches into look-ahead buffer

Special case

	7	6	5	4	3	2	1					output		
he		s	a	i	d	:		H	A	H	A	H	AHA!	(0,0,H)
he	s	a	i	d	:		H	A	H	A	H	A	HH!	(0,0,A)
...e s	a	i	d	:		H	A	H	A	H	A	H	A!	(2,4,H)
...d:	H	A	H	A	H	A	H	A	!					(2,1,!)
... HA	H	A	H	A	H	A	!							

Search buffer	Look-ahead buffer
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Encoding – Pseudo code algorithm

```
while look-ahead buffer is not empty
  go backwards in search buffer to find longest match of the look-ahead buffer

  if match found
    print: (offset from window boundary, length of match, next symbol in look-
    ahead buffer);
    shift window by length+1;
  else
    print: (0, 0, first symbol in look-ahead buffer);
    shift window by 1;
  fi
end while
```


Example (Decoding)

input		7	6	5	4	3	2	1
(0,0,a)								a
(0,0,b)							a	b
(0,0,r)						a	b	r
(3,1,c)				a	b	r	a	c
(2,1,d)		a	b	r	a	c	a	d
(7,4,d)	abrac	a	d	a	b	r	a	d

Decoding – Pseudo code algorithm

```
for each token (offset, length, symbol)
  if offset = 0 then
    print symbol;
  else
    go reverse in previous output by offset characters and copy
    character wise for length symbols;
    print symbol;
  fi
next
```

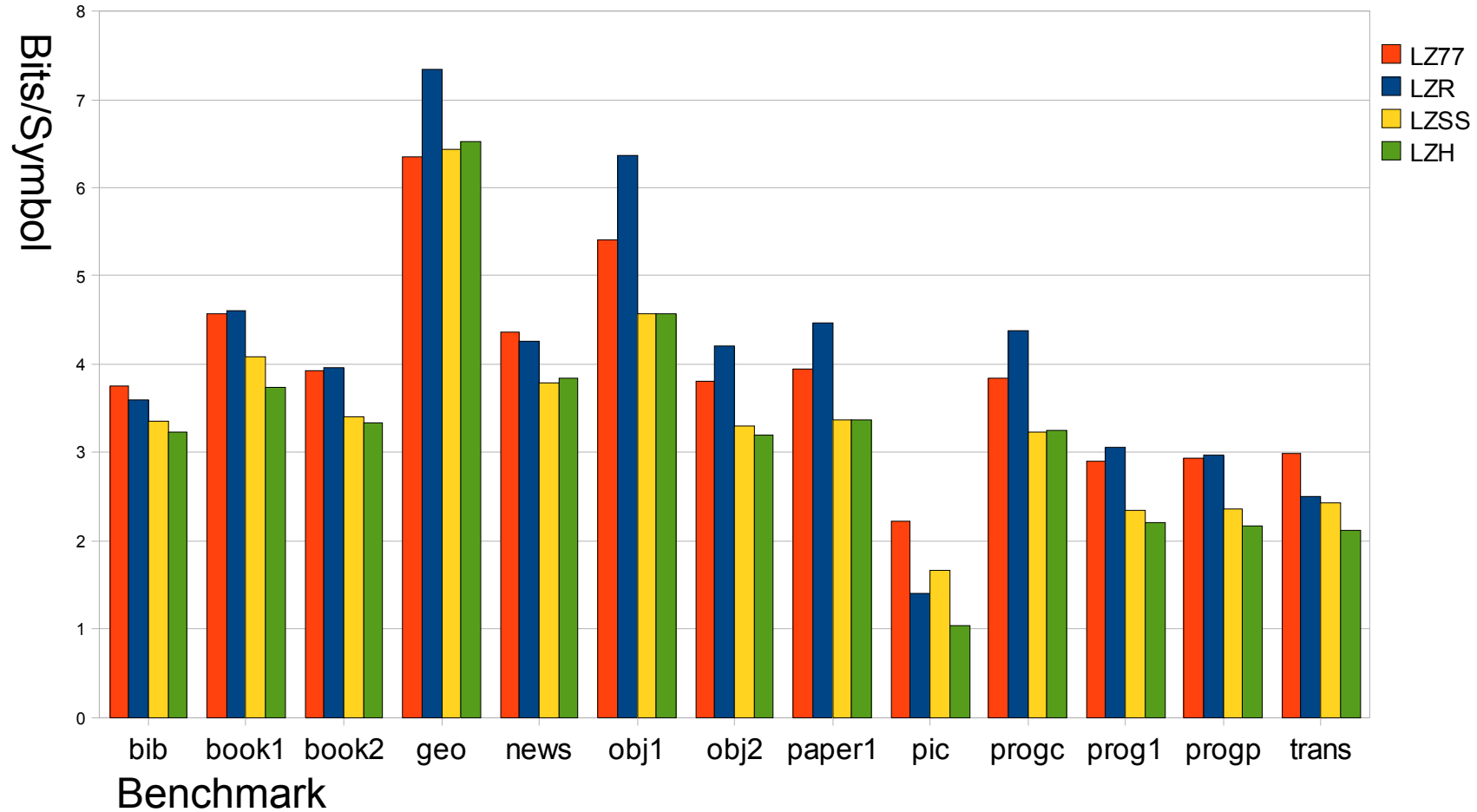
LZ77 is asymmetric, encoding is more difficult than decoding as it needs to find the longest match.

Optimizations

Successors following LZ77 used different optimizations:

- Use variable size offset and length fields in the tuples instead of fixed-length. Better if small offsets and sizes prevail.
 - Don't output a (0,0,x) token when character is not found but instead differentiate using a flag-bit: 0|x or 1|o,l
 - Use better suited data structure (e.g. tree, hash set) for the buffers. This allows faster search and/or larger buffers.
 - Additional Huffman coding of tuples/references.
- > LZSS, LZB, LZH, LZR, LZFG, LZMA, Deflate, ...

Performance



(From Bell/Cleary/Witten: Text Compression)

Applications, Patents

Unlike **LZ78**, **LZ77** has not been patented. This may be a reason why its successors basing on LZ77 are so widely used:

Deflate is a combination of LZSS together with Huffman encoding and uses a window size of 32kB.

This algorithm is open source and used in what is widely known as ZIP compression (although the ZIP format itself is only a container format, like AVI and can be used with several algorithms), and by the formats PNG, TIFF, PDF and many others.

References

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